# Al Classical and Non-deterministic Planning: Model-based Autonomous Behavior

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# Part I

# Non-deterministic Planning

# Part 1: Non-deterministic Planning

- 1 Non-deterministic Planning
- 2 Solution Concepts for FOND Planning
- 3 Solving FOND Planning
  - FOND Planning using Classical Planners
  - FOND Planning via SAT
  - Compact Policies via ASP/SAT
- 4 Conditional Fairness

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# Planning Models: Vanilla Model for Classical Al Planning

- finite and discrete state space S
- a known initial state  $s_0 \in S$
- a set  $S_G \subseteq S$  of goal states
- actions  $A(s) \subseteq A$  applicable in each  $s \in S$
- a deterministic transition function s' = f(a, s) for  $a \in A(s)$
- positive action costs c(a, s)

A solution/plan is seq. of applicable actions  $\pi = a_0, \dots, a_n$  that maps  $s_0$  into  $S_G$ .

Plan is optimal if it minimizes the sum of action costs.

Different models obtained by relaxing assumptions in bold.

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# Planning with non-deterministic actions

What if an action may yield different effect outcomes?

- Slipery floor: you may slip and fall (and maybe hurt yourself).
- Slipery blocksworld: if you stack or unstack a block, it may fall down to the table.
- **Dice rolling:** if you roll a die, it may yield different outcomes: 1,2,3,4,5 or 6.
- **Robot operation:** when using the gripper, it may succeed or fail to pick an object (and may need to retry).



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- Robot operation: when using the gripper, it may succeed or fail to pick an object (and may need to retry).
- Finding parking: when visiting a block you may or may not find parking space (if not, keep going around the block).
- Walking on beam: if you do a step on a beam, you may advance or fall down.
- Walking on corridor: if you do a step you may or may not be at the end of the corridor.



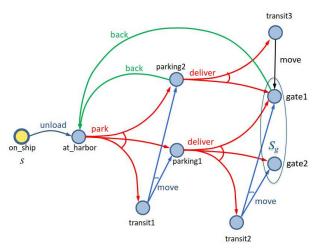
# Example: Harbor Management FOND Problem

Very simple harbor management domain:

- 1 Unload a single item from a ship.
- Park the item in a storage facility.
- 3 Deliver it to gates (to be loaded into tracks).



Storage and gates may be unavailable, but we can always wait and move containers around.



(Example 11.1 in *Acting, Planning, and Learning* Ghallab, Nau, Traverso 2025)

# Planning with Markov Decision Processes

#### Goal MDPs are fully observable, probabilistic state models:

- $\blacksquare$  a state space S
- **2** initial state  $s_0 \in S$
- 3 a set  $G \subseteq S$  of goal states
- 4 actions  $A(s) \subseteq A$  applicable in each state  $s \in S$
- **5** transition probabilities  $P_a(s' \mid s)$  for  $s \in S$  and  $a \in A(s)$  **1**
- 6 action costs c(a,s) > 0
- Solutions are functions (called "policies") mapping states into actions;  $\pi:S\mapsto A$ 
  - $ightharpoonup \pi(s)$  states what action to do in state s
- Optimal solutions minimize expected cost to goal.
- Reward-based MDPs involve rewards instead of costs, and discount factor  $\gamma \in [0,1)$  in place of goals. They underlie theory of RL.  $\ensuremath{\mathfrak{C}}$

# FOND Planning: Fully-observable Non-Deterministic Planning

#### A **FOND state model** is like the "logical" counterpart of Goal MDPs:

- $\blacksquare$  a state space S
- **2** initial state  $s_0 \in S$
- $\mathbf{3}$  a set  $G \subseteq S$  of goal states
- 4 actions  $A(s) \subseteq A$  applicable in each state  $s \in S$
- **5** non-det state transition function F: successors  $s' \in F(a,s)$ ,  $s \in S$ ,  $a \in A(s)$  **3**
- 6 action costs c(a, s) = 1

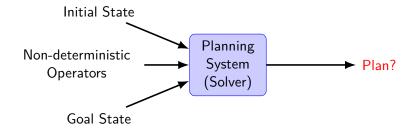
- Main change from Classical Planning: F(a,s) maps to set of possible states (not to one unique state).
  - Nature decides what next state is reached after action a is applied in state s non-determinism.
  - ... but agent will observe the state reached after a is applied.

# FOND Planning: Fully-observable Non-Deterministic Planning

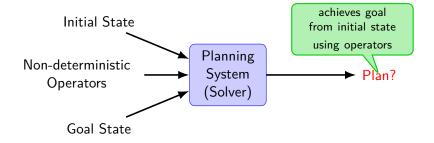
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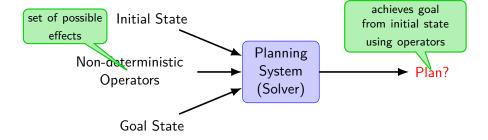
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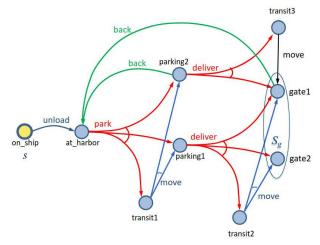
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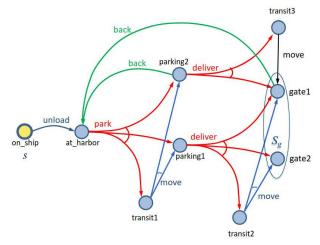


- Is it possible to always deliver the containers to the gates?
- If so, what is the sequence of actions?



(Example 11.1 in Acting, Planning, and Learning Ghallab, Nau, Traverso 2025)

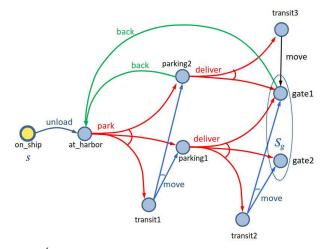
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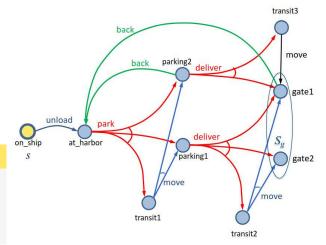
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#### **Policy**

A **policy**  $\pi$  is a partial function from states s into actions a; that is,  $\pi: S \mapsto A$ . (when undefined, agent stops acting)



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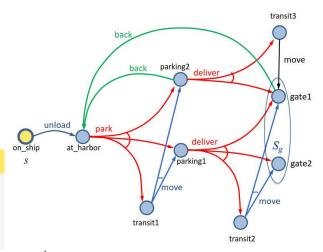
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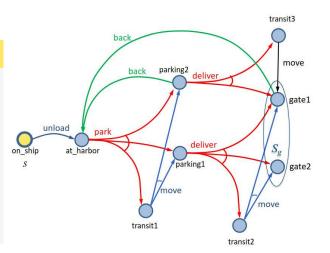
 $\ref{Solution}$  Is there a "good" policy  $\pi$ ?



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# Example: Does $\pi_1$ solve the task?

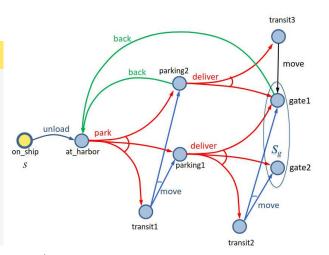
S	$\pi_1(s)$
on_ship	unload
at_harbor	park
parking1	deliver
parking2	back
transit1	move
transit2	move
transit3	move



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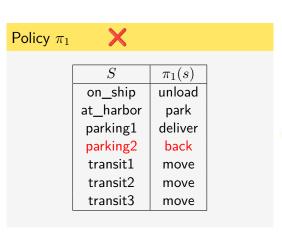
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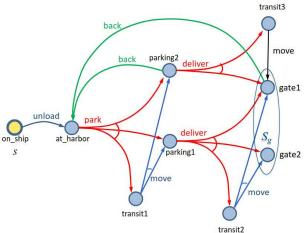
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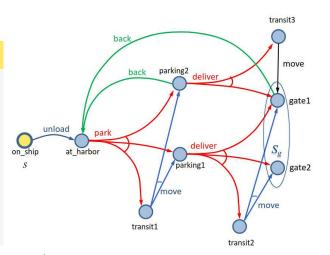




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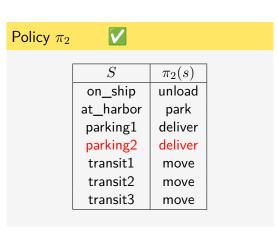
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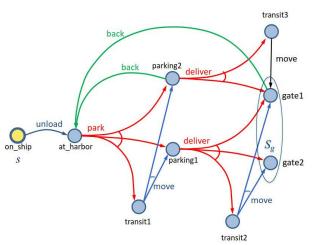
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on_ship	unload
at_harbor	park
parking1	deliver
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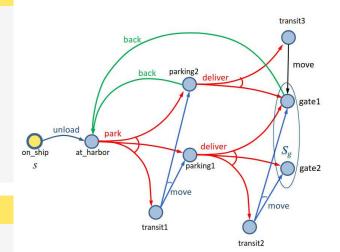
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# Example: Which one is better?

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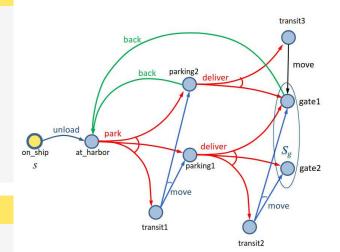


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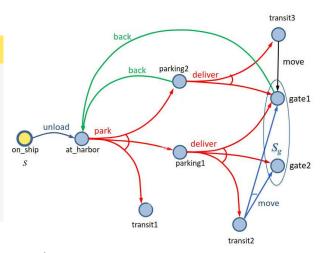
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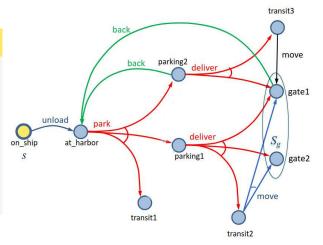
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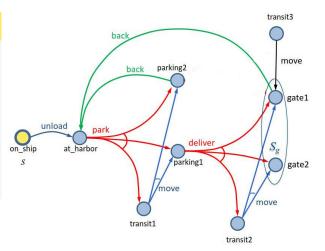
But could  $\pi_2$  succeed (sometimes)?



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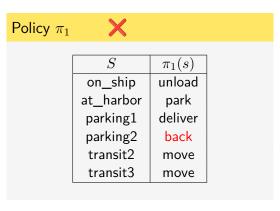
# Example: What if parking2 is not connected to gates?

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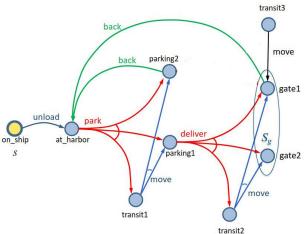


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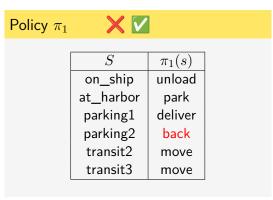


Storage parking1 may never be available!



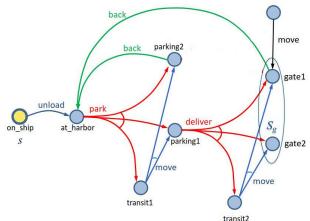
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# Example: What if parking2 is not connected to gates?



Storage parking1 may never be available!

But, what if we know parking1 would eventually becomes available?

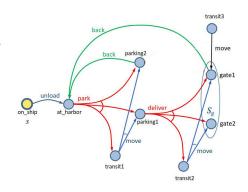


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transit3

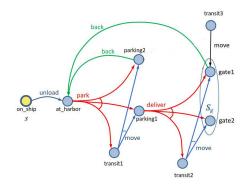
# So, some lessons...

- Classical plans as sequences of actions are not enough to solve FOND problems.
- We need to use a policy that maps states into actions.
  - ► More like "programs" with conditionals and loops!
- Some (bad) policies are better than others.
- Some policies may achieve the goal, but not always.
- Some policies will achieve the goal if environment is not too adversarial — not unfair.



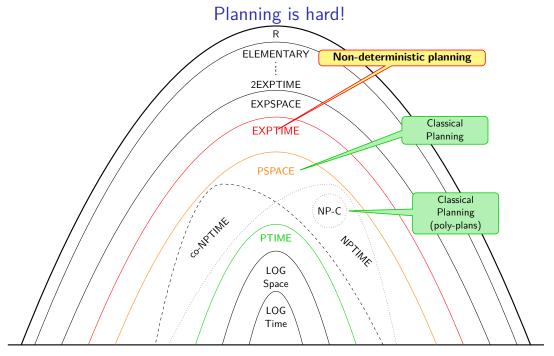
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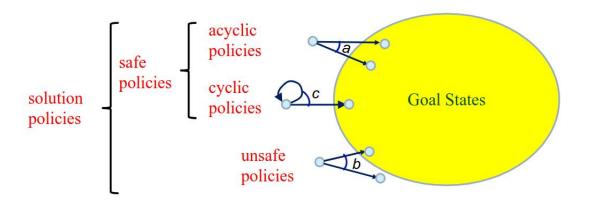


This seems way more complex planning!





#### Kinds of Solution Policies



Acting, Planning, and Learning Ghallab, Nau, Traverso 2025

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Running policy  $\pi$  from state s yields trajectories runs:

- $\pi$ -trajectories  $s_0, \ldots, s_n$ , such that  $s_{i+1} \in F(a_i, s_i)$ ,  $a_i = \pi(s_i)$ , for  $i \in [0, n-1]$ .
- $\pi$ -trajectory **maximal** if 1)  $s_n$  is goal state, 2)  $\pi(s_n) = \bot$ , or 3)  $n = \infty$  ( $\pi$  is infinite)

## **FOND Planning Solution Concepts**

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  - Always a possibility to reach the goal.

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  - ▶ All executions are guaranteed to reach the goal (in a known bounded number of actions!).
  - Plans may have conditionals (but no loops!)
- 3  $\pi$  is strong cyclic solution if for each state s reachable from  $s_0$  with a  $\pi$ -trajectory, there is a  $\pi$ -trajectory from s to goal.
  - Always a possibility to reach the goal.
  - ► Goal will be achieved if environment is not "adversarial"

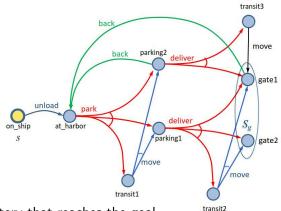
Running policy  $\pi$  from state s yields trajectories runs:

- $\pi$ -trajectories  $s_0, \ldots, s_n$ , such that  $s_{i+1} \in F(a_i, s_i)$ ,  $a_i = \pi(s_i)$ , for  $i \in [0, n-1]$ .
- $\pi$ -trajectory **maximal** if 1)  $s_n$  is goal state, 2)  $\pi(s_n) = \bot$ , or 3)  $n = \infty$  ( $\pi$  is infinite)

- **1**  $\pi$  is a **weak solution** if there is a  $\pi$ -trajectory from  $s_0$  that reaches goal.
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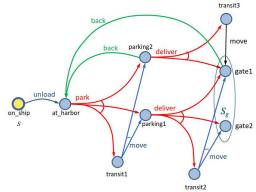
### Weak Plans

S	$\pi_1(s)$
on_ship	unload
at_harbor	park
parking1	deliver
parking2	back
transit2	move
transit3	move



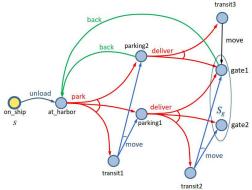
- ✓ Policy  $\pi$  is a weak plan as there is a trajectory that reaches the goal.
  - ► {on\_ship}, {at\_harbor}, {parking1}, {gate1}
- $\star$  But  $\pi$  is *not* a strong plan.
  - ▶ {on\_ship}, {at\_harbor}, {parking2}, {at\_harbor}, {parking2}, {at\_harbor}, ...

	( )
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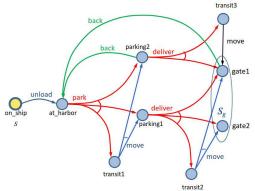


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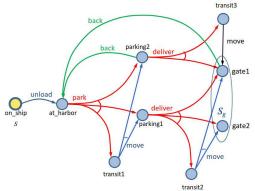
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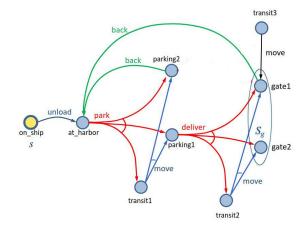
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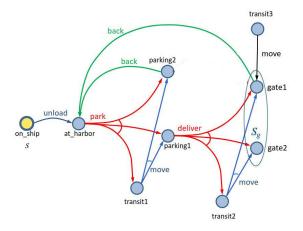
Policy  $\pi$  is **strong cyclic solution** if for each state s reachable from  $s_0$  with a  $\pi$ -trajectory, there is a  $\pi$ -trajectory from s to goal.

- Yes!, policy never "loses" the possibility to get the goal
- But, it may loop "forever" in some states.
- We can make  $\pi$  strong by changing it to  $\pi_1(\text{parking2}) = \text{deliver}$ .

? Is there a strong plan?



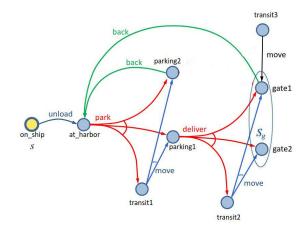
Is there a strong plan? No!



**?** Is there a strong plan? **No!** Best we can do is:

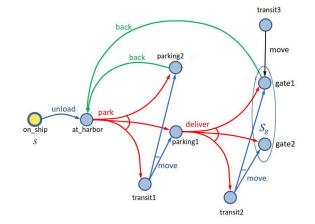
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**?** When will this policy reach the goal?



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When will this policy reach the goal?
When executed in "fair" environments!

### Fairness Environments

# Non-determinism behavior under fairness assumption

A strong cyclic policy eventually reaches the goal in every **fair** trajectory.



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What type of environments?



#### Fairness Environments

## Non-determinism behavior under fairness assumption

A strong cyclic policy eventually reaches the goal in every **fair** trajectory.

- **?** What type of environments?
  - Where each effect listed has indeed non-zero probability.
  - Re-trying is an effective strategy.
    - rolling a die until it shows a 6.
    - driving around the block until a parking space is available.
    - pour into cup until full.



### Fairness Environments

- Classical sequential plans are not enough to solve FOND problems.
  - ▶ We need more flexible behavior description (controlller) for agents
- We use **policies** mapping states into actions.
  - Allow conditional and loops.

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## Question

How can we compute these plans with loops? How to compute strong-cyclic plans policies?

# Part 1: Non-deterministic Planning

- 1 Non-deterministic Planning
- 2 Solution Concepts for FOND Planning
- 3 Solving FOND Planning
  - FOND Planning using Classical Planners
  - FOND Planning via SAT
  - Compact Policies via ASP/SAT
- 4 Conditional Fairness

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### Non-determinism in PDDI

- Non-deterministic effects added to PDDI for the 5th IPC in 2006.
- Action effect can have a one-of effect:

```
(oneof e1 e2 ... en)
```

To support uncertainty track in IPC-5.

#### 5th International Planning Competition: Non-deterministic Track Call For Participation

#### Blai Bonet Departamento de Computación

Universidad Simón Bolívar Caracas, Venezuela bonet@ldc.usb.ve

#### Abstract

The 5th International Planning Competition will be colocated with ICAPS-06. This IPC edition will contain a track on nondeterministic and probabilistic planning as the continuation of the probabilistic track at IPC-4. The non-deterministic track will evaluate systems for conformant, non-deterministic and probabilistic planning under different criteria. This document describes the general goals of the track, the planning tasks to be addressed, the representation language and the evaluation methodology.

#### Introduction

The 5th International Planning Competition (IPC-5) will be colocated with the 16th International Conference on Automated Planning and Scheduling, ICAPS-06, to be held in The English Lake District, UK, during June 6-10, 2006. The IPC is a biannual event where planning systems are evalu-

#### Robert Givan

Electrical & Computer Engineering Purdue University West Lafavette, IN 47907 givan@ecn.purdue.edu

tracks that will cover the areas of non-deterministic conformant planning, non-deterministic planning (i.e. conditional planning with full observability), and probabilistic planning (i.e. conditional probabilistic planning with full observabil-

As done in the classical track of IPC, we believe that planners that offer different guarantees on the quality of their solutions should be evaluated differently; otherwise the comparisons are not meaningful. Hence, planners within each group will be further categorized by the guarantees they provide, as much as possible given the number of participants.

The rest of this document is organized as follows. Sect. 2 gives a brief background on the different planning tasks included in the competition as well as the form of the solutions. Sect. 3 presents the extensions and restrictions upon the PPDDL language to be used. Sect. 4 focuses on the evaluation aspects of the competition, mainly how different

```
(:action unstack
    :parameters (?b1 ?b2 - block)
    :precondition (and (not (= ?b1 ?b2)) (emptyhand) (clear ?b1) (on ?b1 ?b2))
    :effect (oneof
     (and (holding ?b1) (clear ?b2) (not (emptyhand)) (not (clear ?b1)) (not (on ?b1 ?b2)))
      (and (clear ?b2) (on-table ?b1) (not (on ?b1 ?b2)))))
            ;; second effect: fail to grab; ?b1 ends on table
```

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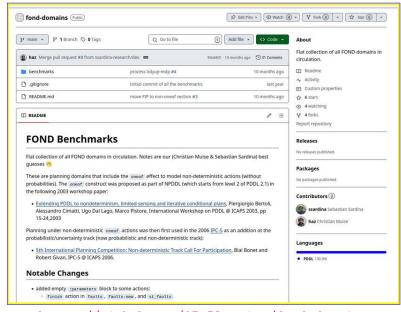
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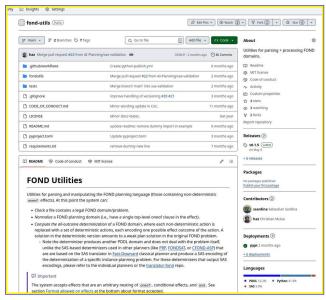
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# Al-Planning/fond-domains @ GH: Benchmark for FOND



https://github.com/AI-Planning/fond-domains

# Al-Planning/fond-utils @ GH: Utilities for FOND



https://github.com/AI-Planning/fond-utils

# FOND Planning using Classical Planners

One of the most effective ways to solve FOND planning problems is to use **classical planners**! Weird...?

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One of the most effective ways to solve FOND planning problems is to use **classical planners**! Weird...?

They all use a **deterministic relaxation** of the FOND problem:

#### All-outcome determinization

**Deterministic relaxation**  $P_D$  of FOND P obtained by substituing **non-det** actions a with effects  $\{e_1, \ldots, e_n\}$  by **deterministic** actions  $a^1, \ldots, a^n$ , where  $a^i$ 's effect is  $e_i$ , for  $i \in [1, n]$ .

- ullet  $P_D$  is a deterministic classical planning problem.
- Under reasonable assumptions,  $P_D$  is polynomially larger than P.
- There are tools to do the determinization: https://github.com/AI-Planning/fond-utils

# Week and Online Solutions for FOND Planning

# $\maltese$ Weak (open loop) solution for P

### From any classical plan $\rho$ for $P_D$ :

- If  $\rho$  generates trajectory  $s_0, \ldots, s_N$  in  $P_D$ , set  $\pi(s_i) = a$  if  $\rho_i \in a$ .
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Reach goal by interacting with FOND "system" that returns **observation**  $s' \in F(a, s)$ :

- **I** From current state s, initially  $s_0$ , compute plan  $\rho = \rho_1, \ldots, \rho_N$  for  $P_D[s]$ .
- **2** Execute **prefix**  $a_1, \ldots, a_i$  for  $\rho_i \in a_i$  until state  $s_i$  **observed** is goal or **different** than state  $s_i'$  **predicted** in  $P_D$ .
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- $\stackrel{\bullet}{\triangleright}$  Properties: If no dead-end states reachable in P, under mild assumptions, goal state eventually reached. Else, method is **incomplete**.

### PRP: Strong Cyclic Policies using Classical Planners

More powerful off-line method, can compute **strong cyclic policies**:

#### PRP: Planning for Relevant Policies (Muise, McIllraith, Beck ICAPS'12)

- **I** Run **simulated on-line** method not just from  $s_0$  but from every possible sucessor s' of a (simulated) **observed** state s; i.e.,  $s' \in F(a, s)$  for a executed in s.
- **2** If state  $s' \in F(a,s)$  is reached from which **no classical plan** for  $P_D(s)$ ; **remove** a from A(s), and start all over again.
- **8** Keep policy to  $\pi(s) = a$  where deterministic version  $a_i$  is head of **shortest classical prefix** found from s to goal.

#### Properties:

Method is sound and complete: returns strong cyclic policy if one exists.



- More **scalable** than other methods as it uses **classical planners**.
- Can be made more efficient by generalizing plans using regression.
- Struggles if there are many "risky" nondeterminism leading to dead-ends.

#### Consider the following situation:

- **1** Goal is  $G = \{g\}$ .
- **2** Classical plan  $\rho = a_1, \dots, a_n$  optimally achieves G from state  $s_0$  in  $P_D$ .
- **3** So,  $\rho$  yields trajectory  $s_0, s_1, \ldots, s_n$  in  $P_D$  such that  $g \in s_n$ .
  - ▶ The last action of  $\rho$  has  $g \in Add(a_n)$   $a_n$  achieves the goal.
- **4** The precondition of  $a_n$  is  $Pre(a_n) = \{p, q\}$ .
  - Clearly,  $p, q \in s_{n-1}$   $a_n$ 's precondition hold just before the goal.

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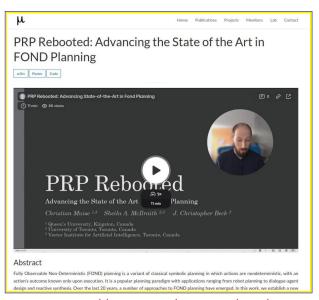
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If  $Add(a_{n-1}) = \{p\}$  and  $Pre(a_{n-1}) = \{w\}$ , what states s' can we set  $\pi(s') = a_{n-1}$ ?

#### PRP Rebooted: AAAI'24



https://mulab.ai/project/pr2/

#### Shortcomings of Classical Planners for FOND

PRP scales wellas it uses **classical planners** + **regression**. However:

- Codebase is highly **sophisticated**; thousands of lines.
- Uses a lot of tricks: regression, dead-end detection, regression, loop closing, strong-cyclic check, etc.
- Struggle from "risky nondeterminism", where previous search choices need to be thrown and restarted.
  - ▶ non-deterministic actions whose other effects will eventually lead to dead-ends.
- May output very large policies no guarantees of "compactness".
- Unable to handle mixed fairness environments.
  - some actions are fair, others are unfair.

#### Shortcomings of Classical Planners for FOND

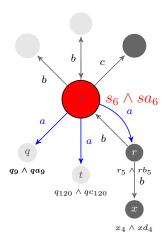
PRP scales wellas it uses classical planners + regression. However:

- Codebase is highly sophisticated; thousands of lines.
- Uses a lot of tricks: regression, dead-end detection, regression, loop closing, strong-cyclic check, etc.
- Struggle from "risky nondeterminism", where previous search choices need to be thrown and restarted.
  - ▶ non-deterministic actions whose other effects will eventually lead to dead-ends.
- May output very large policies no guarantees of "compactness".
- Unable to handle **mixed fairness** environments.
  - some actions are fair, others are unfair.
- What can we do about these issues? Can we get a simpler, declarative solver for FOND?

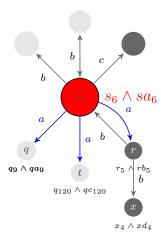
# Recall Theory C(P, n) for Classical Problem $P = \langle F, A, I, G \rangle$

- Init:  $p_0$  for  $p \in I$ ,  $\neg q_0$  for  $q \in F$  and  $q \notin I$
- Goal:  $p_n$  for  $p \in G$
- Actions: For  $i = 0, 1, \dots, n-1$ , and each action  $a \in A$ 
  - $ightharpoonup a_i \supset p_i \text{ for } p \in Prec(a)$
  - $ightharpoonup a_i\supset p_{i+1}$  for each  $p\in Add(a)$
  - $ightharpoonup a_i \supset \neg p_{i+1}$  for each  $p \in Del(a)$
- **Persistence:** For  $i=0,\ldots,n-1$ , and each atom  $p\in F$ , where  $O(p^+)$  and  $O(p^-)$  stand for the actions that add and delete p resp.
  - $\triangleright p_i \land \land_{a \in O(p^-)} \neg a_i \supset p_{i+1}$
  - $ightharpoonup \neg p_i \wedge \wedge_{a \in O(p^+)} \neg a_i \supset \neg p_{i+1}$
- **Seriality:** For each  $i=0,\ldots,n-1$ , if  $a\neq a'$ ,  $\neg(a_i \wedge a_i')$

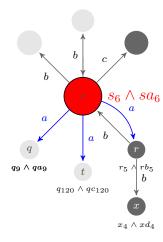
Key idea: label each state with action and distance to goal.



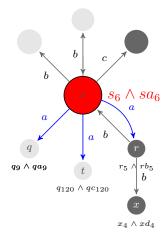
- - **Key idea:** label each state with action and distance to goal.
  - Variables of SAT encoding (i is not time index!)
    - $ightharpoonup s_i$ : min "distance" from s to goal in policy is at most i
    - $ightharpoonup sa_i$ :  $s_i$  and  $\pi(s) = a$



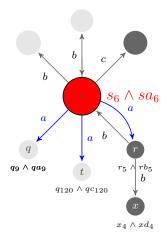
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  - Formulas C(M); here  $M = \mathcal{S}(P)$  and max = |S| 1:
    - **1**  $s_{max}$  for initial state  $s_I$ ; max dist I to goal of length  $\leq max$
    - 2  $s_0$  for  $s \in S_G$  and  $\neg s_0$  for  $s \notin S_G$



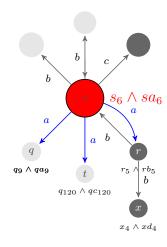
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    - $s_i \supset \bigvee_{a \in A(s)} sa_i$  ; choose action in s, preserve distance



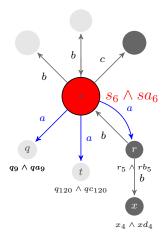
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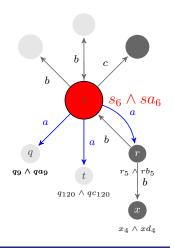
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    - 7  $sa_{max}\supset s'_{max}$  ; if  $\pi(s)=a$ , all  $s'\in f(a,s)$ , must reach goal
    - 8  $sa_{max} \supset \neg sa'_{max}$ ; if  $\pi(s) = a$ , then  $\pi(s) \neq a'$ ,  $a \neq a'$ .



- **Key idea:** label each state with action and distance to goal.
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#### Theorem

- **1** Model M has a **strong-cyclic policy**  $\pi$  iff C(M) is satisfiable.
- 2 If  $\sigma$  satisfies C(M),  $\pi(s)=a$  for  $sa_{max}$  true in  $\sigma$  is a **strong-cyclic policy** that solves M

#### Too large encoding: Towards Compact Polocies

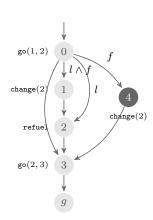
- Encodings are **exhaustive**, all states s represented! \*
- (Geffner & Geffner 2018) proposed an encoding in SAT computing compact policies.
  - of course, not in worst case
- Can also be adjusted to compute strong policies.
- Can also handle dual FOND: fair and unfair actions!
- (Yadav & Sardina 2023): alternative encoding in a Answer Set Programming (ASP):
  - ► More compact exploits ASP first-order language.
  - ▶ More readable uses a more declarative style.
  - ► Integrates regression ideas from PRP.
  - Exploits ASP technology.

## Compact Controllers via ASP (Yadav & Sardina 2023)

**Key idea:** devise a finite state controller with n states - (Geffner & Geffner 2018)

Encoding in ASP for FOND problem  $P = \langle A, I, G \rangle$ :

- atom(P): for each predicate  $P \in A$ .
- action(A): for each action A ∈ A. In addition, to define an action's precondition and effects we use the following terms:
  - ▶ prec(A, P): atom P is in precondition of action A.
  - effect(A, E): associates an action with its E-th effect (one per oneoff effect).
  - ▶ add(A, E, P): E-th effect of action A adds atom P.
  - ▶ del(A, E, P): E-th effect of action A deletes atom P.
- init(P): predicate  $P \in I$  is true in the initial state.
- goal(P): predicate  $P \in G$  is in the goal condition.



#### Define Controllers States and Transitions

#### Solver to decide:

- policy(S, A): action A executed in controller state S.
- 2 next(S1, E, S2): S2 is the next controler state if the E-th effect of prescribed action in S1 ocurrs.

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```
state(0..k). % states of the controller
[policy(S, A): action(A)] = 1:- state(S), S != k.
[next(S1, E, S2): state(S2)] = 1 :- policy(S1, A), effect(A, E).
```

- **1** Defines controller k+1 states. State k is goal state.
- **2** Select one action per controller state (except goal state k).
- 3 Defines a transition for each action's effect to a next controller state.

#### Define Controllers States and Transitions

```
1 holds(S, P) :- policy(S, A), prec(A, P).
2 holds(S1, P) :-
3    next(S1, E, S2), holds(S2, P), policy(S1, A), not add(A, E, P).
4 -holds(S2, P) :- next(S1, E, S2), policy(S1, A), del(A, E, P).
5 -holds(0, P) :- atom(P), not init(P).
6 holds(k, P) :- goal(P).
```

- Preconditions must hold where action is prescribed.
- 2
- **3** Regression: P must have been true in the previous controller state.
- 4 Progression: P must be false next if action deleted it.
- 5 Initial state negative atoms.
- 6 What must be true at goal controller state k

### Define Solution Concept: Strong Cyclic

```
reachableG(S):- state(S), S = k.
reachableG(S):- next(S, _, S1), reachableG(S1).
reachableG(S), state(S).
```

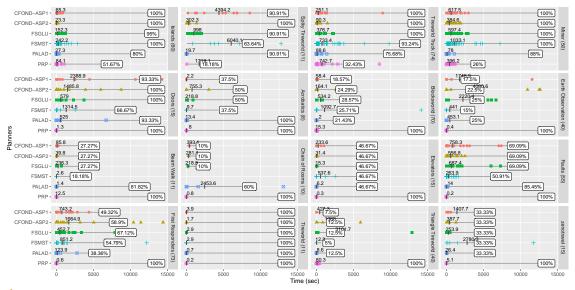
- Goal controller state is reachable from itself.
- 2 Transitive clousure: Any (previous) controller state connected to a controller state that reaches the goal state, also reaches the controller goal state.
- **3 Constraint:** No controller state does not reach the goal state.

#### Full FOND-ASP Code

```
1 state(0..k). % states of the controller
2 {policy(S, A): action(A)} = 1:- state(S), S != k.
3 {next(S1, E, S2): state(S2)} = 1 :- policy(S1, A), effect(A, E).
  holds(S, P) := policy(S, A), prec(A, P).
  holds(S1, P) :-
     next(S1, E, S2), holds(S2, P), policy(S1, A), not add(A, E, P).
8 -holds(S2, P) :- next(S1, E, S2), policy(S1, A), del(A, E, P).
9 -holds(0, P) :- atom(P), not init(P).
no holds(k, P) := goal(P).
11
reachableG(S):- state(S), S = k.
reachableG(S):- next(S, _, S1), reachableG(S1).
14 :- not reachableG(S), state(S).
```

☆ If a model is returned, controller defined in predicates policy/2 and next/3.

#### Experimental Results vs. PRP and FOND-SAT



Better in risky non-determinism domains — first five. PRP better in the rest.

### Recap SAT/ASP for FOND Planning

- Declarative elegant solver for FOND planning problems via SAT or ASP.
- Compact controllers: finite state controller with k+1 states.
- Increase the size when no solution found, and repeat.
- Faster than classical planning based approaches in domains with meaningful non-determinism ("risky").
- Can incorporate domain control knowledge (e.g., "do not executre a after b").
- Still struggles with large domains with "easy" non-determinism.

#### Part 1: Non-deterministic Planning

- 1 Non-deterministic Planning
- 2 Solution Concepts for FOND Planning
- 3 Solving FOND Planning
  - FOND Planning using Classical Planners
  - FOND Planning via SAT
  - Compact Policies via ASP/SAT
- 4 Conditional Fairness

#### Part 1: Non-deterministic Planning

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Consider an robot in a corridor:

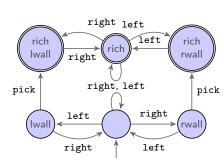


- Robot can move *left* or right (up to the walls). Unknown size of steps, but  $\geq 1$
- A price is at some of the end of the corridor.
- Robot doesn't know its cell, but can sense if there is a wall on left/right after moving.
- **?** Can the robot get the money? How to model the setting?

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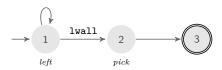
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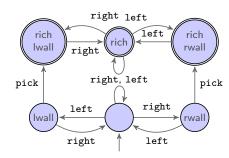


#### Consider an robot in a corridor:



Would this controller work?

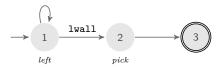




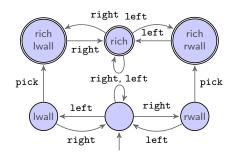
Consider an robot in a corridor:



Would this controller work? YES!



**Strong-cyclic policy:** Retrying *left* works!

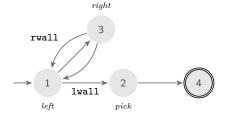


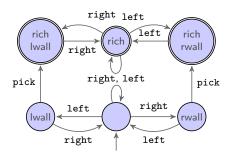
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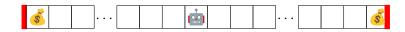
What about this one?



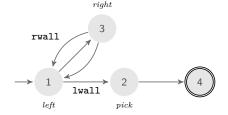


# Can the robot get the money?

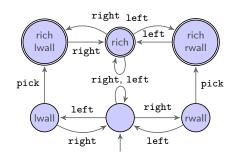
Consider an robot in a corridor:



What about this one? NO!



How come? It is also a **strong-cyclic policy!** States where rich true are always reachable.. *left* action done infinitely many times in initial state



### Conditional Fairness (Rodriguez et al. 2021)

- Standard fairness assumption is not enough:
  - trying *left* is not sufficient!
  - ightharpoonup must not move right while trying...



- We need conditional fairness: left is fair as long as right is not executed.
  - Same for right: fair provided left is not executed.
- Standard FOND planners cannot handle this: they assume that all actions are fair.
- (Rodriguez et al. 2021)'s FOND<sup>+</sup> in ASP can handle:
  - Strong-cyclic policies with conditional fairness.
  - Mixed fairness: some actions are fair, others not.



(Best Paper Award ICAPS'21)

### FOND<sup>+</sup>

Let's generalize FOND:

#### FOND<sup>+</sup> Problem

A FOND<sup>+</sup> problem  $P_c = \langle P, C \rangle$  is a FOND problem P extended with a set C of **(conditional) fairness assumptions** of the form  $A_i/B_i$ ,  $i=1,\ldots,n$  and where each  $A_i$  is a set of **non-deterministic actions** in P, and each  $B_i$  is a set of actions in P disjoint from  $A_i$ .

<u>Meaning of  $A/B \in C$ </u>: If a state trajectory contains infinite occurrences of actions  $a \in A$  in a state s, and *finite* occurrences of actions from B, then s must be immediately followed by each  $s' \in F(\pi(s), s)$  an infinite number of times.

if left is executed infinitely many times in s, but right is executed, say, 10 times, then eventually we will see the left wall.

### FOND Solutions as FOND<sup>+</sup> Solutions

#### FOND<sup>+</sup> Problem

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Strong and strong cyclic planning all have solutions defined by the implicit fairness assumptions particular to each one of them.

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Strong and strong cyclic planning all have solutions defined by the implicit fairness assumptions particular to each one of them.

#### Theorem

The **strong-cyclic solutions** of a FOND problem P are the solutions of the FOND<sup>+</sup> problem  $P_c = \langle P, \{A/\emptyset\} \rangle$ , where A is the set of all the non-deterministic actions in P.

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Strong and strong cyclic planning all have solutions defined by the implicit fairness assumptions particular to each one of them.

#### $\mathsf{Theorem}$

The **strong-cyclic solutions** of a FOND problem P are the solutions of the FOND<sup>+</sup> problem  $P_c = \langle P, \{A/\emptyset\} \rangle$ , where A is the set of all the non-deterministic actions in P.

#### Theorem

The **strong solutions** of a FOND problem P are the solutions of the FOND<sup>+</sup> problem  $P_c = \langle P, \emptyset \rangle$ .

### FOND<sup>+</sup>-ASP: An ASP-based Planner

```
% policy, edges, and connectedness
                                                                               STATE(S)
   \{ pi(S,A) : ACTION(A) \} = 1 :- STATE(S), not GOAL(S).
2
                                                                               INITIAL(S)
   successor(S,T) :- pi(S,A), TRANSITION(S,A,T).
                                                                               GOAL(S)
                                                                               ACTION(A)
                                                                               TRANSITION(S,A,T)
   connected(S,T) :- successor(S,T).
                                                                               ASET(A,I)
   connected(S,T) :- connected(S,X), successor(X,T), S != X.
                                                                               BSET(B,I)
   blocked(S,T) :- STATE(S), STATE(T), not connected(S,T).
   blocked(S,T) :- connected(S,T), terminate(S).
   blocked(S,T) :- connected(S,T), terminate(T).
   blocked(S,T) :- connected(S,T),
12
                    blocked(X.T): successor(S.X), connected(X.T).
   fair(S) := pi(S,A), ASET(I,A), blocked(X,S) : pi(X,B), BSET(I,B), not blocked(S,X).
   % terminating states
   terminate(S) :- GOAL(S).
   terminate(S): - fair(S), successor(S,T), terminate(T).
   terminate(S) :- not fair(S), successor(S,_), terminate(T) : successor(S,T)
20
   % reachable states must terminate
   :- reachable(S), not terminate(S).
   reachable(S) :- INITIAL(S).
   reachable(S) :- reachable(X), not GOAL(X), successor(X,S).
```

# FOND<sup>+</sup>-ASP: Graphical Intuition...

figure of a transition system, with two states looping, the first selects action A and the second B. draw successors of each..

### FOND<sup>+</sup>-ASP: Solution Policy

```
1  % policy, edges, and connectedness
2  { pi(S,A) : ACTION(A) } = 1 :- STATE(S), not GOAL(S).
3  successor(S,T) :- pi(S,A), TRANSITION(S,A,T).
4
5  % reachable states must terminate
6 :- reachable(S), not terminate(S).
7  reachable(S) :- INITIAL(S).
8  reachable(S) :- reachable(X), not GOAL(X), successor(X,S).
```

- 2 Select an action per domain state.
- 3 Edges are transitions of the action selected.

### FOND<sup>+</sup>-ASP: Solution Policy

```
1  % policy, edges, and connectedness
2  { pi(S,A) : ACTION(A) } = 1 :- STATE(S), not GOAL(S).
3  successor(S,T) :- pi(S,A), TRANSITION(S,A,T).
4
5  % reachable states must terminate
6 :- reachable(S), not terminate(S).
7  reachable(S) :- INITIAL(S).
8  reachable(S) :- reachable(X), not GOAL(X), successor(X,S).
```

- 2 Select an action per domain state.
- **3** Edges are transitions of the action selected.
- 6 **Constraint:** every reachable state via the policy needs to eventually terminate.
- 7-8 Define reachable states via the policy.

### FOND<sup>+</sup>-ASP: State Termination

Defines when a state will eventually lead to termination and not get "sucked" in a loop..

- 2 If the state is a goal state.
- 3 If state will behave **fairly** (wrt effects of prescribed action) and one successor state will terminate.
- 4 If state may *not* behave **fairly**, and all successors will terminate.

### FOND<sup>+</sup>-ASP: Fairness

```
connected(S,T) :- successor(S,T).
   connected(S,T) := connected(S,X), successor(X,T), S != X.
3
   % terminating states
   terminate(S) :- GOAL(S).
   terminate(S) :- fair(S), successor(S,T), terminate(T).
   terminate(S) :- not fair(S), successor(S,_),
                    terminate(T) : successor(S,T).
   fair(S) := pi(S,A), ASET(I,A),
10
              blocked(X,S): pi(X,B), BSET(I,B), not blocked(S,X).
```

- 1-2 States connected by the policy.
- 4-7 Every path from s to T will terminate somewhere.
- 10 Fair if any loop that includes actions in a fairness pair A/B (e.g., left and right), will terminate somewhere else.

## FOND<sup>+</sup>-ASP: Strong Cyclic

#### **Theorem**

The **strong-cyclic solutions** of a FOND problem P are the solutions of the FOND<sup>+</sup> problem  $P_c = \langle P, \{A/\emptyset\} \rangle$ , where A is the set of all the non-deterministic actions in P.



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always false

fair(S) :- pi(S,A), ASET(I,A), always false

blocked(X,S) : pi(X,B), BSET(I,B), not blocked(S,X).
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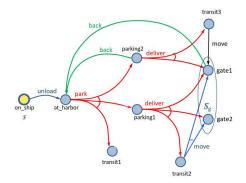
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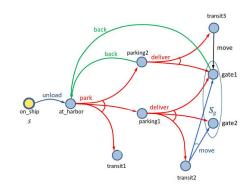
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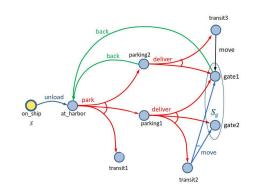
- We tested FOND<sup>+</sup>-ASP experimentally:
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  - Just add oneof in effects!
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- FOND<sup>+</sup> and domains with "qualitative" numbers?
  - e.g., distance to the wall

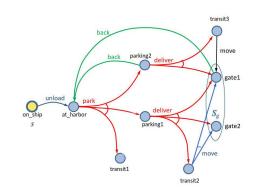


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Qualitative Numeric Planning (QNP)



# Que vimos? ••

- Busqueda as a general problem solving method:
  - Representación: state model (a graph!).
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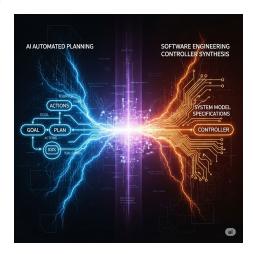
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  - Planning graphs.
- **3 FOND Planning**: Non-determinism
  - ► Non-deterministic state models (no probabilities!)
  - ▶ PDDL with one-of effects + Policies.
  - Solution concepts: weak, strong, strong-cyclic.
  - Fairness assumption on environment.
  - Computing policies.



# Al Planning and Control Synthesis in SE 🤝

- What if we want to plan for more complex goals?
  - **Elevator controller:** every passenger floor requests needs to be *eventually* fulfilled, but **never** have more than 10 passengers on board.
- Event-driven systems: some events cannot be planned/controlled (e.g., user aborts transaction)
- Infinite behavior: continuous operation, never stop.
  - What are the goals if we never finish? Infinite games vs. finite games
- Compositional planning/synthesis: software components described separately
  - Plan on different PDDLs and the combine.



# LaFHIS - Laboratory on Fundamentals and Tools for Software Engineering





The Tools and Foundations for Software Engineering Lab

What we do The Lab News Contact

# R&D Augmentation

We help organisations solve difficult problems by applying state of the art automated software engineering methods, techniques and tools. We support our partners in bootstrapping their R&D activities, designing strategies, identifying key technologies and collaboratively developing solutions.

We incorporate, combine and adapt state of the art techniques from program analysis, program repair, program understanding domain, specific programming languages, and model-based software engineering as needed to address the specific contexts and bottlenecks that our partners have.



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